

# Characterizing the inscribability of polytopes using slack matrices

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# The problem

Given a polytope, does there exist a combinatorially equivalent\* polytope with all vertices on the sphere?

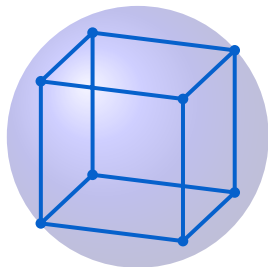
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\*The face lattices of the two polytopes are isomorphic

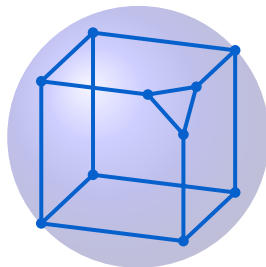
# The problem

Given a polytope, does there exist a combinatorially equivalent\* polytope with all vertices on the sphere?

Yes!  $\Rightarrow$  Inscriptible



No!  $\Rightarrow$  Non-inscribable



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\*The face lattices of the two polytopes are isomorphic

## Solving a nonlinear system

**Theorem.** [Firsching, 2017] A  $d$ -polytope  $P$  with  $n$  vertices is inscribable if and only if the following system has a solution

$$\chi(i_1, \dots, i_r) \det(\bar{p}_{i_1}, \dots, \bar{p}_{i_r}) > 0 \quad \text{for all } i_1, \dots, i_r \in \binom{[n]}{r}$$
$$\|p_i\|^2 = 1 \quad \text{for } i = 1, \dots, n$$

where  $r = d + 1$ ,  $\bar{p}_i = \begin{bmatrix} p_i \\ 1 \end{bmatrix}$ ,  $[n] = \{1, \dots, n\}$ , and  $\chi : [n]^r \rightarrow \{-1, 0, 1\}$  is the associated chirotope of  $P$

**Note:** The system has  $nd$  variables,  $\binom{n}{r}$  inequalities of degree  $r$ , and  $n$  equalities of degree 2

Suppose  $P$  is a  $d$ -polytope such that

- $P = \text{conv}\{v_1, \dots, v_n\}$
- $P = \{x \in \mathbb{R}^d : 1 - h_j^\top x \geq 0, j = 1, \dots, m\}$

Let  $V = [v_1 \cdots v_n] \in \mathbb{R}^{d \times n}$  and  $H = [h_1 \cdots h_m] \in \mathbb{R}^{d \times m}$

A slack matrix  $S \in \mathbb{R}^{n \times m}$  of  $P$  is

$$S = \begin{bmatrix} \mathbf{1}_n & V^\top \end{bmatrix} \begin{bmatrix} \mathbf{1}_m^\top \\ -H \end{bmatrix} = \mathbf{1}_{n \times m} - V^\top H$$

**Note:**  $S_{ij} = 1 - h_j^\top v_i$  is facet inequality  $j$  evaluated on vertex  $i$

**Theorem.** A  $d$ -polytope  $P$  is inscribable if and only if there exists

$$X = \begin{bmatrix} 1 & \mathbb{1}_n^\top & \mathbb{1}_m^\top \\ \mathbb{1}_n & A & S \\ \mathbb{1}_m & S^\top & B \end{bmatrix} \succcurlyeq 0$$

such that

- $\text{rank}(X) = d + 1$
- $\text{diag}(A) = \text{const.}$
- $S \geq 0$
- $S$  has the same support as slack matrices of  $P$

# Characterizing inscribability using slack matrices

Let  $I^Z$  be the index set of zeros in slack matrices

Consider the following optimization problem

$$\begin{aligned} \min_X \quad & \text{rank}(X) \\ \text{s.t.} \quad & X = \begin{bmatrix} \mathbf{1} & \mathbb{1}_n^\top & \mathbb{1}_m^\top \\ \mathbb{1}_n & A & S \\ \mathbb{1}_m & S^\top & B \end{bmatrix} \succcurlyeq 0 \\ & S_{ij} = 0, \text{ if } (i, j) \in I^Z \\ & S_{ij} > 0, \text{ if } (i, j) \notin I^Z \\ & A_{ii} = 2, i = 1, \dots, n. \end{aligned}$$

**Note:** The minimum of this problem is no less than  $d + 1$

**Question:** How to solve this?

## **An SDP formulation**



## An SDP formulation

Instead of solving the min-rank problem directly, we consider the following SDP problem

$$\begin{aligned} \min_X \quad & \text{tr}(X) - \sum_{(i,j) \notin I^Z} \lambda_{ij} S_{ij} \\ \text{s.t.} \quad & X = \begin{bmatrix} 1 & \mathbb{1}_n^\top & \mathbb{1}_m^\top \\ \mathbb{1}_n & A & S \\ \mathbb{1}_m & S^\top & B \end{bmatrix} \succeq 0 \\ & S_{ij} = 0, \text{ if } (i,j) \in I^Z \\ & A_{ii} = 2, i = 1, \dots, n \end{aligned} \tag{P}$$

where  $\lambda_{ij}$  are some positive weights

# Dual problem and duality gap

The dual problem of (P) is

$$\begin{aligned} \max_{u,v} \quad & m + n + \sum_{\substack{1 \leq i \leq n \\ 1 \leq j \leq m}} M_{ij} - \sum_{i=1}^n u_i + 1 \\ \text{s.t.} \quad & \begin{bmatrix} I_n + \text{diag}(u) & \frac{1}{2}M \\ \frac{1}{2}M^T & I_m \end{bmatrix} \succcurlyeq 0 \end{aligned} \tag{D}$$

where

$$\text{diag}(u) = \begin{bmatrix} u_1 & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & u_n \end{bmatrix} \quad \text{and} \quad M_{ij} = \begin{cases} -\lambda_{ij}, & \text{if } (i,j) \notin I^Z \\ v_k \text{ for some } k, & \text{if } (i,j) \in I^Z \end{cases}$$

**Note:** The Slater's condition holds so the duality gap is zero

# Works for inscribable polytopes?

For an inscribable polytope  $P$

- Find an inscription of  $P$  and a set of weights  $\lambda_{ij}$  such that the duality gap is zero evaluated at the inscription
- Then, solving the SDP with the  $\lambda_{ij}$  above will give an inscription of  $P$

## Example: $n$ -gons

For  $n$ -gons, we have  $n = m \geq 3$  and  $d = 2$

Problem (P) evaluated at the inscription

$$v_i = \left[ \cos \frac{2(i-1)\pi}{n} \quad \sin \frac{2(i-1)\pi}{n} \right]^\top, i = 1, \dots, n$$

$$h_j = \frac{1}{\cos \frac{\pi}{n}} \left[ \cos \frac{(2j-1)\pi}{n} \quad \sin \frac{(2j-1)\pi}{n} \right]^\top, j = 1, \dots, n$$

with

$$\lambda_{ij} = \frac{2}{n \cos^2 \frac{\pi}{n}}, \forall (i, j) \notin I^Z$$

has

$$f_p^* = f_d^* = 2n - n \tan^2 \frac{\pi}{n} + 1$$

**Therefore, the SDP formulation solves the inscribability problem for  $n$ -gons**

**For  $d$ -simplices,  $d$ -cubes, and  $d$ -crosspolytopes, the SDP formulation solves the inscribability problem**

In particular, solving (P) with the following weights gives an inscription

$$(d\text{-simplex}) \quad \lambda_{ij} = \frac{2d^2}{d+1}, \quad \forall (i,j) \notin I^z$$

$$(d\text{-cube}) \quad \lambda_{ij} = d2^{1-d}, \quad \forall (i,j) \notin I^z$$

$$(d\text{-crosspolytope}) \quad \lambda_{ij} = 1, \quad \forall (i,j) \notin I^z$$

# Algorithms

# Algorithms for the min-rank problem

- Solving the SDP formulation
- Alternating projection: project between the rank  $d + 1$  cone and the feasible set  $\Omega$  of the min-rank problem
- Simplified alternating projection: replace the projection onto  $\Omega$  with forcing  $X_k$  to have correct constants on correct positions

# Numerical experiments

- Test set: inscribable simplicial 4-polytopes with  $5 \leq n \leq 10$  vertices
- Starting point for AP and simplified AP is the solution of the SDP
- Results comparing accuracy:

Alg \ n	5	6	7	8	9	10
SDP	100%	100%	80%	84%	87%	59%
AP (simplified)	100%	100%	80%	86%	89%	61%
AP	100%	100%	100%	95%	97%	94%

- Results comparing runtime (s):

Alg \ n	5	6	7	8	9	10
SDP	0.4	0.5	0.6	0.9	1.3	1.5
AP (simplified)	0.4	0.5	1.1	1.5	2.7	7.16
AP	0.4	0.5	15.1	86.7	122.0	330.2



In this research, we

- Characterized the inscribability problem of polytopes as a min-rank optimization problem based on slack matrices
- Proposed an SDP formulation to approximate the min-rank problem and proved it works for certain classes of polytopes
- Provided three algorithms with different accuracy and efficiency for the inscribability problem

Thank you